

Laver Tables: from Set Theory to Braid Theory

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Joint work with Patrick DEHORNOY

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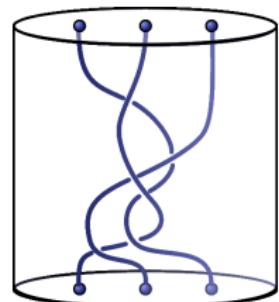
Topology Symposium, Tohoku University, July 29, 2014

A_3	1	2	3	4	5	6	7	8
1	2	4	6	8	2	4	6	8
2	3	4	7	8	3	4	7	8
3	4	8	4	8	4	8	4	8
4	5	6	7	8	5	6	7	8
5	6	8	6	8	6	8	6	8
6	7	8	7	8	7	8	7	8
7	8	8	8	8	8	8	8	8
8	1	2	3	4	5	6	7	8

$0, 1, 2, 3, \dots;$

$\aleph_0, \aleph_1, \aleph_2, \dots;$

\aleph_ω, \dots



Part 1

A Laver table is...

Basic definitions

A **shelf** (= self-distributive structure)

is a set S with an operation \triangleright satisfying

$$a \triangleright (b \triangleright c) = (a \triangleright b) \triangleright (a \triangleright c) \quad (\text{SD})$$

Example: group G , $f \triangleright g = fgf^{-1}$.

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is a free shelf generated by a single element γ .

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Theorem (Laver, '95): properties (SD) and (Init) uniquely define \triangleright .

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Theorem (Laver, '95): properties (SD) and (Init) uniquely define \triangleright .

⚠ False for $\{1, 2, 3, \dots, q\}$ with $q \neq 2^n$.

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$$\begin{array}{c} \gamma \\ \gamma \triangleright \gamma \end{array}$$

$$\begin{array}{c} (\gamma \triangleright \gamma) \triangleright \gamma \\ ((\gamma \triangleright \gamma) \triangleright \gamma) \triangleright \gamma \\ \dots \end{array}$$

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$$\gamma = 1$$

$$(\gamma \triangleright \gamma) \triangleright \gamma = 3$$

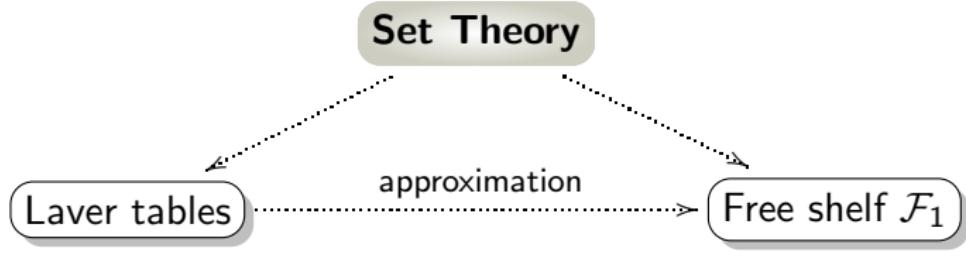
$$\gamma \triangleright \gamma = 2$$

$$((\gamma \triangleright \gamma) \triangleright \gamma) \triangleright \gamma = 4$$

...

Laver tables in Set Theory

Richard Laver



Laver tables in Set Theory: details

“Super-infinite” sets

Finite \iff every **self-embedding** is bijective.

Infinite \iff admits a non-bijective self-embedding.

Example: \mathbb{N} is infinite ($n \mapsto n + 1$ is a non-bijective self-embedding),

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Axiom I3

V_λ (a certain limit rank) is super-infinite.

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V_λ (a certain limit rank) is super-infinite.

⚠ I3 can neither be proved nor refuted in Zermelo-Fraenkel system.

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Self-embeddings

Set $S \rightsquigarrow \text{Emb}(S) := \{ f : S \hookrightarrow S \} \rightsquigarrow \text{a shelf } (\text{Emb}(S), \triangleright)$

$$f \triangleright g = \begin{cases} fgf^{-1} & \text{on the image } \text{Im}(f) \text{ of } f, \\ \text{Id} & \text{on the complement of } \text{Im}(f). \end{cases}$$

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I3 \Rightarrow $\diamondsuit f_0$ generates a sub-shelf $F \subseteq \text{Emb}(V_\lambda)$, with $F \cong \mathcal{F}_1$;

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 - ✿ F has quotients of size $2^n \rightsquigarrow$ **Laver tables!**

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 - ✿ F has quotients of size $2^n \rightsquigarrow$ **Laver tables!**
 - ✿ $\varprojlim_{n \in \mathbb{N}} A_n \supseteq \mathcal{F}_1 \rightsquigarrow A_n$ are **finite approximations** of \mathcal{F}_1

Going beyond Set Theory?

Elementary definition

$A_n = (\{1, 2, 3, \dots, 2^n\}, \triangleright)$ satisfying

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Elementary properties

✿ A projective system of shelves:

$$\begin{aligned} p_n : A_n &\longrightarrow A_{n-1}, \\ a &\longmapsto a \pmod{2^{n-1}}. \end{aligned}$$

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✿ Periodic rows:

$p \triangleright 1$	$<$	$p \triangleright 2$	$<$	\cdots	$<$	$p \triangleright 2^r$	\dots	periodic repetition ...
$= p + 1$						$= 2^n$		

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$\pi_n(p) := 2^r$ is the **period** of p in A_n .

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$$A_n \cong \mathcal{F}_1 / (\cdots ((\gamma \triangleright \gamma) \triangleright \gamma) \cdots) \triangleright \gamma = \gamma$$

$$1 \leftrightarrow \gamma$$

$$2 \leftrightarrow \gamma \triangleright \gamma$$

$$3 \leftrightarrow (\gamma \triangleright \gamma) \triangleright \gamma$$

$$4 \leftrightarrow ((\gamma \triangleright \gamma) \triangleright \gamma) \triangleright \gamma \quad \dots$$

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- Some “nice” rows and columns

A_n	1	2	3	.	2^{n-1}	.	2^n	π_n
					
2^{n-1}	$2^{n-1}+1$	$2^{n-1}+2$	$2^{n-1}+3$...	2^n	...	2^n	2^{n-1}
				
2^n-3	2^n-2	2^n	2^n-2	...	2^n	...	2^n	2
2^n-2	2^n-1	2^n	2^n-1	...	2^n	...	2^n	2
2^n-1	2^n	2^n	2^n	...	2^n	...	2^n	1
2^n	1	2	3	...	2^{n-1}	...	2^n	2^n

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2^n-3	2^n-2	2^n	2^n-2	...	2^n	...	2^n	2
2^n-2	2^n-1	2^n	2^n-1	...	2^n	...	2^n	2
2^n-1	2^n	2^n	2^n	...	2^n	...	2^n	1
2^n	1	2	3	...	2^{n-1}	...	2^n	2^n

⚠ No closed formulas for $p \triangleright q$, nor for $\pi_n(p)$.

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- ❖ $\pi_n(1) \underset{n \rightarrow \infty}{\rightarrow} \infty$.

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- ❖ $\pi_n(1) \underset{n \rightarrow \infty}{\rightarrow} \infty.$
- ❖ $\pi_n(1) \leq \pi_n(2).$

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- ⚠ Theorems under Axiom I3!

A_0	1
1	1

A_2	1	2	3	4
1	2	4	2	4
2	3	4	3	4
3	4	4	4	4
4	1	2	3	4

A_1	1	2
1	2	2
2	1	2

A_3	1	2	3	4	5	6	7	8
1	2	4	6	8	2	4	6	8
2	3	4	7	8	3	4	7	8
3	4	8	4	8	4	8	4	8
4	5	6	7	8	5	6	7	8
5	6	8	6	8	6	8	6	8
6	7	8	7	8	7	8	7	8
7	8	8	8	8	8	8	8	8
8	1	2	3	4	5	6	7	8

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1	1

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1	2	4	2	4
2	3	4	3	4
3	4	4	4	4
4	1	2	3	4

A_1	1	2
1	2	2
2	1	2

A_3	1	2	3	4	5	6	7	8	
1	2	4	6	8	2	4	6	8	$\pi_3(1) = 4$
2	3	4	7	8	3	4	7	8	$\pi_3(2) = 4$
3	4	8	4	8	4	8	4	8	$\pi_3(3) = 2$
4	5	6	7	8	5	6	7	8	$\pi_3(4) = 4$
5	6	8	6	8	6	8	6	8	$\pi_3(5) = 2$
6	7	8	7	8	7	8	7	8	$\pi_3(6) = 2$
7	8	8	8	8	8	8	8	8	$\pi_3(7) = 1$
8	1	2	3	4	5	6	7	8	$\pi_3(8) = 8$

A_0	1
1	1

A_2	1	2	3	4
1	2	4	2	4
2	3	4	3	4
3	4	4	4	4
4	1	2	3	4

A_1	1	2
1	2	2
2	1	2

A_3	1	2	3	4	5	6	7	8	
1	2	4	6	8	2	4	6	8	$\pi_3(1) = 4$
2	3	4	7	8	3	4	7	8	$\pi_3(2) = 4$
3	4	8	4	8	4	8	4	8	$\pi_3(3) = 2$
4	5	6	7	8	5	6	7	8	$\pi_3(4) = 4$
5	6	8	6	8	6	8	6	8	$\pi_3(5) = 2$
6	7	8	7	8	7	8	7	8	$\pi_3(6) = 2$
7	8	8	8	8	8	8	8	8	$\pi_3(7) = 1$
8	1	2	3	4	5	6	7	8	$\pi_3(8) = 8$

A_4	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
1	2	12	14	16	2	12	14	16	2	12	14	16	2	12	14	16
2	3	12	15	16	3	12	15	16	3	12	15	16	3	12	15	16
3	4	8	12	16	4	8	12	16	4	8	12	16	4	8	12	16
4	5	6	7	8	13	14	15	16	5	6	7	8	13	14	15	16
5	6	8	14	16	6	8	14	16	6	8	14	16	6	8	14	16
6	7	8	15	16	7	8	15	16	7	8	15	16	7	8	15	16
7	8	16	8	16	8	16	8	16	8	16	8	16	8	16	8	16
8	9	10	11	12	13	14	15	16	9	10	11	12	13	14	15	16
9	10	12	14	16	10	12	14	16	10	12	14	16	10	12	14	16
10	11	12	15	16	11	12	15	16	11	12	15	16	11	12	15	16
11	12	16	12	16	12	16	12	16	12	16	12	16	12	16	12	16
12	13	14	15	16	13	14	15	16	13	14	15	16	13	14	15	16
13	14	16	14	16	14	16	14	16	14	16	14	16	14	16	14	16
14	15	16	15	16	15	16	15	16	15	16	15	16	15	16	15	16
15	16	16	16	16	16	16	16	16	16	16	16	16	16	16	16	16
16	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16

A_4	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
1	2	12	14	16	2	12	14	16	2	12	14	16	2	12	14	16
2	3	12	15	16	3	12	15	16	3	12	15	16	3	12	15	16
3	4	8	12	16	4	8	12	16	4	8	12	16	4	8	12	16
4	5	6	7	8	13	14	15	16	5	6	7	8	13	14	15	16
5	6	8	14	16	6	8	14	16	6	8	14	16	6	8	14	16
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8	9	10	11	12	13	14	15	16	9	10	11	12	13	14	15	16
9	10	12	14	16	10	12	14	16	10	12	14	16	10	12	14	16
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11	12	16	12	16	12	16	12	16	12	16	12	16	12	16	12	16
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13	14	16	14	16	14	16	14	16	14	16	14	16	14	16	14	16
14	15	16	15	16	15	16	15	16	15	16	15	16	15	16	15	16
15	16	16	16	16	16	16	16	16	16	16	16	16	16	16	16	16
16	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16

A_4	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
1	2	12	14	16	2	12	14	16	2	12	14	16	2	12	14	16
2	3	12	15	16	3	12	15	16	3	12	15	16	3	12	15	16
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15	16	16	16	16	16	16	16	16	16	16	16	16	16	16	16	16
16	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16

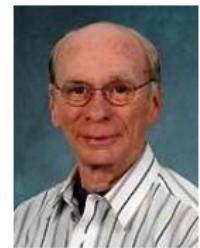
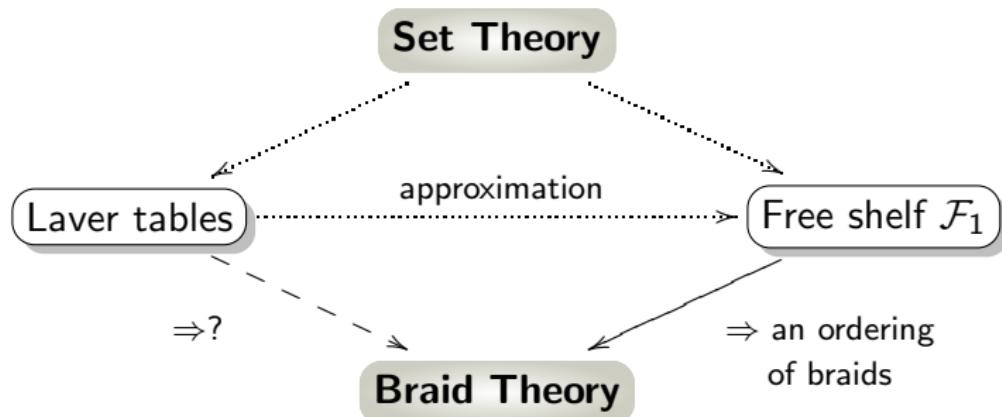
⚠ Rich combinatorics.

Part 2

*Dreams: braid and knot invariants
based on Laver tables*

Laver tables in Topology

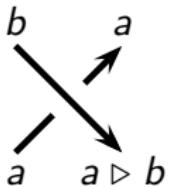
Richard Laver



Patrick Dehornoy

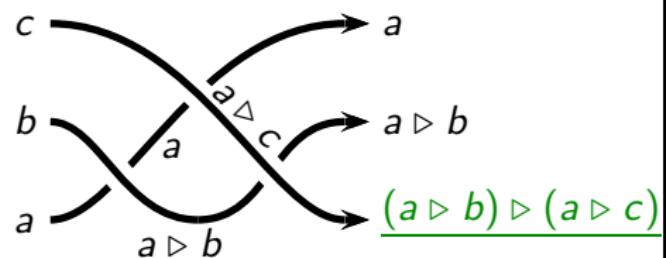
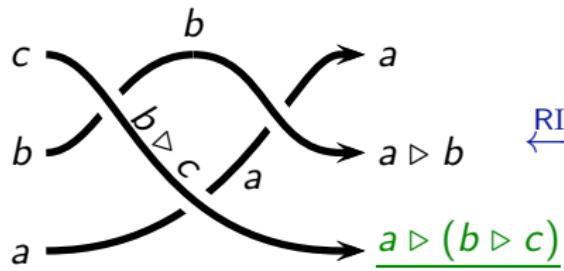
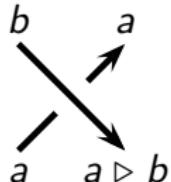
Shelf colorings

Colorings
by (S, \triangleright) :



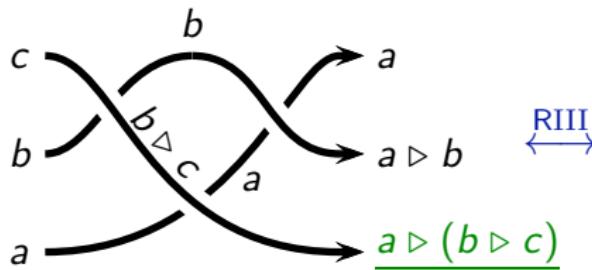
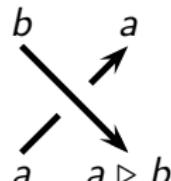
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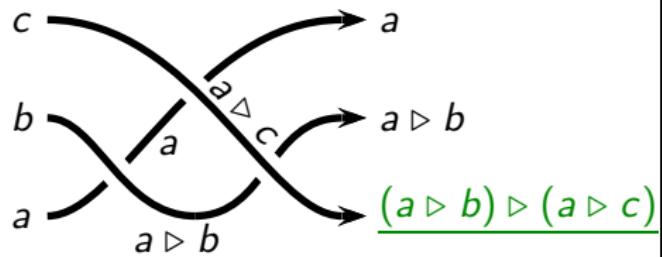


Shelf colorings

Colorings
by (S, \triangleright) :



$\xleftarrow{\text{RIII}}$

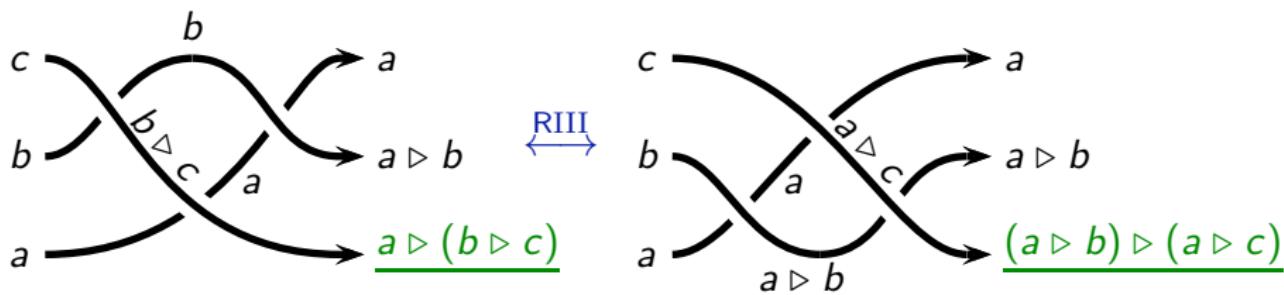
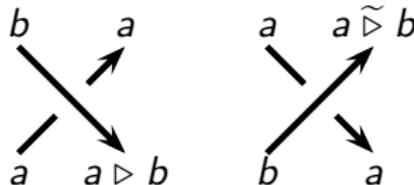


$$\text{RIII} \leftrightarrow a \triangleright (b \triangleright c) = (a \triangleright b) \triangleright (a \triangleright c) \quad (\text{SD})$$

positive braid invariants $\stackrel{\text{colorings}}{\leadsto}$ shelf

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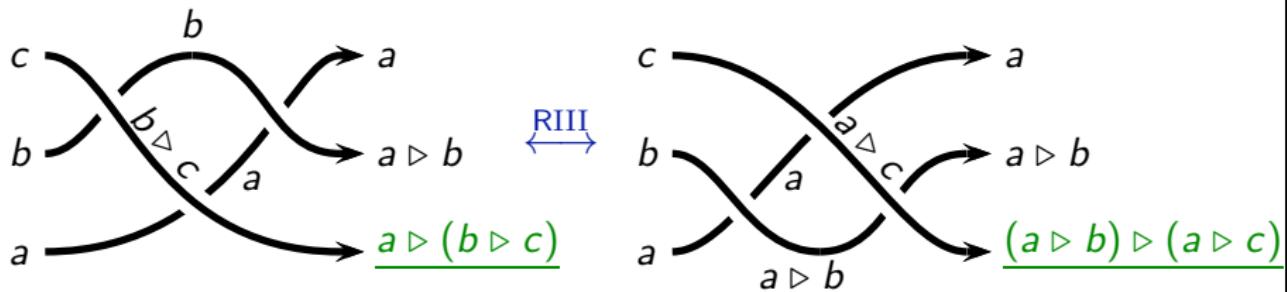
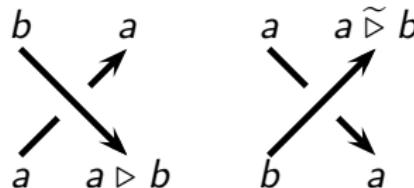
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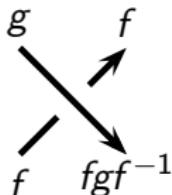


$$\begin{array}{lcl} \text{RIII} & \leftrightarrow & a \triangleright (b \triangleright c) = (a \triangleright b) \triangleright (a \triangleright c) \quad (\text{SD}) \\ \text{RII} & \leftrightarrow & a \tilde{\triangleright} (a \triangleright b) = b = a \triangleright (a \tilde{\triangleright} b) \quad (\text{Inv}) \end{array}$$

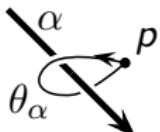
Rack

braid invariants $\overset{\text{colorings}}{\sim}$ rack

Shelf colorings

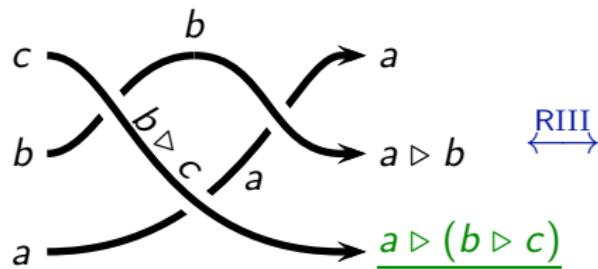


Wirtinger presentation:

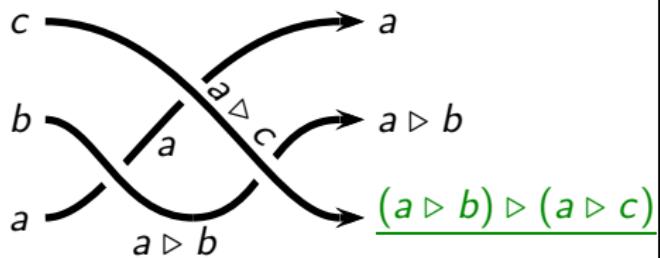


colorings by G

$$\text{Rep}(\pi_1((\mathbb{R}^2 \times [0, 1]) \setminus \beta), G)$$



$\xleftarrow{\text{RIII}}$



RIII	\leftrightarrow	$a \triangleright (b \triangleright c) = (a \triangleright b) \triangleright (a \triangleright c)$	(SD)	}
RII	\leftrightarrow	$a \widetilde{\triangleright} (a \triangleright b) = b = a \triangleright (a \widetilde{\triangleright} b)$	(Inv)	

Rack

braid invariants $\stackrel{\text{colorings}}{\leadsto}$ rack

Example: Group $G \rightsquigarrow$ a rack
 $(G, f \triangleright g = f g f^{-1}, f \widetilde{\triangleright} g = f^{-1} g f)$.

\mathcal{F}_1 -colorings for arbitrary braids?

positive braid invariants $\stackrel{\text{colorings}}{\leadsto} \mathcal{F}_1$

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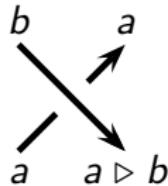
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(S, \triangleright) is a rack \iff the **color propagation map** is **bijective**.

$$\begin{aligned}\sigma : S \times S &\longrightarrow S \times S, \\ (a, b) &\longmapsto (a \triangleright b, a).\end{aligned}$$



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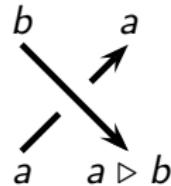
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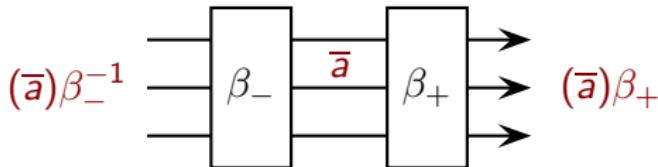
For \mathcal{F}_1 , the map σ is only **injective** \implies partially invertible.

\mathcal{F}_1 -colorings for arbitrary braids?

Normal form for braids: $\boxed{\beta = \beta_- \beta_+}$, β_- negative, β_+ positive.

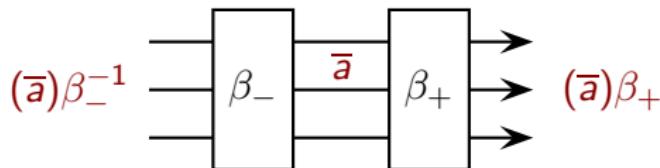
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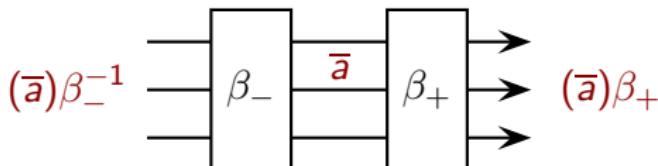


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✿ $\forall k$ -braids β, β' , \exists a common propagable color vector \bar{a} .

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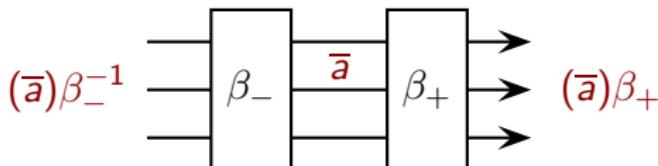


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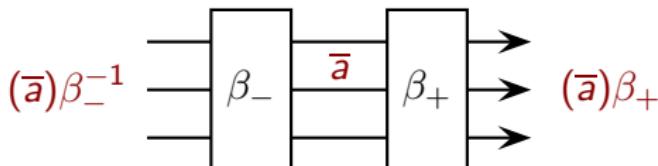
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- ✿ Applications: ✓ efficient algorithms for distinguishing braids,
✓ geometry of closed braids, etc.

A_n -colorings for arbitrary braids?

positive braid invariants $\stackrel{\text{colorings}}{\leadsto}$ Laver tables

Question: What about arbitrary braid invariants?

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$$\sigma(2^n - 1, b) = ((2^n - 1) \triangleright b, 2^n - 1) = (2^n, 2^n - 1).$$

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Part 3

*Reality: 2- and 3-cocycles
for Laver tables*

Shelf colorings revisited

Aim: Add flexibility to coloring invariants.

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Method: enrich colorings with [weights](#)

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Rack cohomology (Fenn-Rourke-Sanderson, '95)

Shelf $(S, \triangleright) \rightsquigarrow$ complex $(\text{Hom}(S^{\times k}, \mathbb{Z}), d_{\text{R}}^k) \rightsquigarrow H_{\text{R}}^k(S)$

$$(d_{\text{R}}^k f)(a_1, \dots, a_{k+1}) = \sum_{i=1}^{k+1} (-1)^{i-1} (f(a_1, \dots, a_{i-1}, a_i \triangleright a_{i+1}, \dots, a_i \triangleright a_{k+1}) - f(a_1, \dots, a_{i-1}, a_{i+1}, \dots, a_{k+1})).$$

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2-cocycles: maps $\phi : S \times S \rightarrow \mathbb{Z}$ satisfying

$$\boxed{\phi(a, c) + \phi(a \triangleright b, a \triangleright c) = \phi(b, c) + \phi(a, b \triangleright c)}$$

Cocycle invariants

Fix a 2-cocycle $\phi : S \times S \rightarrow \mathbb{Z}$:

$$\phi(a, c) + \phi(a \triangleright b, a \triangleright c) = \phi(b, c) + \phi(a, b \triangleright c)$$

ϕ -weight (Carter-Jelsovsky-Kamada-Langford-Saito, '99):

$$\boxed{\begin{array}{ccc} S\text{-colored diagram} & \longmapsto & \sum_{\substack{b \\ a}} \phi(a, b) \end{array}}$$

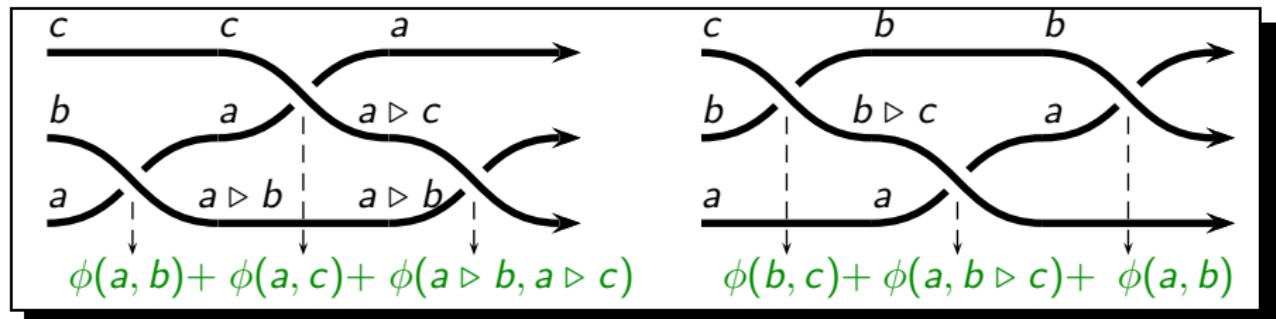
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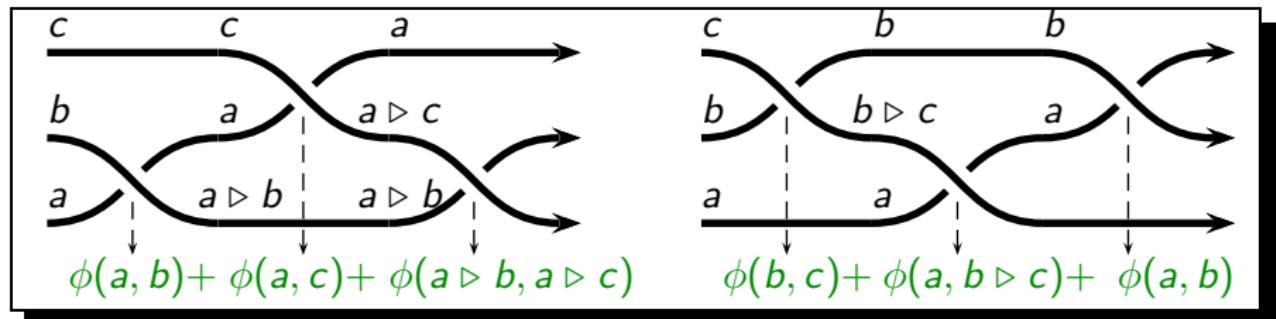
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positive braid invariants

colorings &
 \sim
weights

shelf & 2-cocycle

Shadow cocycle invariants

Fix a 3-cocycle $\psi : S \times S \times S \rightarrow \mathbb{Z}$:

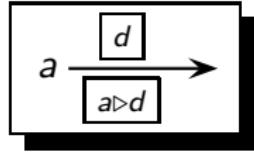
$$\begin{aligned} \psi(a, b, c \triangleright d) + \psi(a, c, d) + \psi(a \triangleright b, a \triangleright c, a \triangleright d) = \\ \psi(b, c, d) + \psi(a, b \triangleright c, b \triangleright d) + \psi(a, b, d) \end{aligned}$$

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Shadow colorings:



ψ -weight:

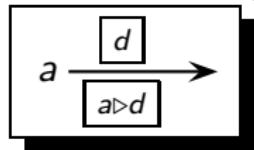
$$\text{S-colored diagram} \quad \mapsto \quad \sum_{\substack{b \triangleright d \\ a \not\sim}} \psi(a, b, d)$$

Shadow cocycle invariants

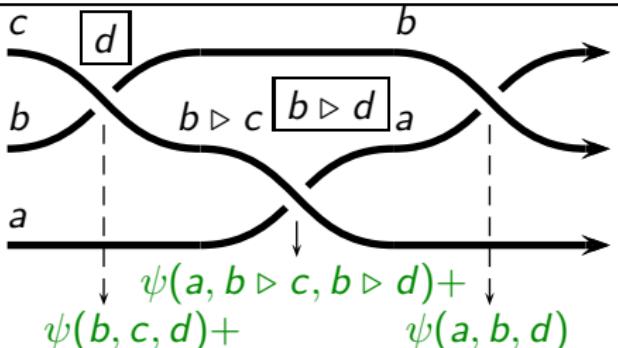
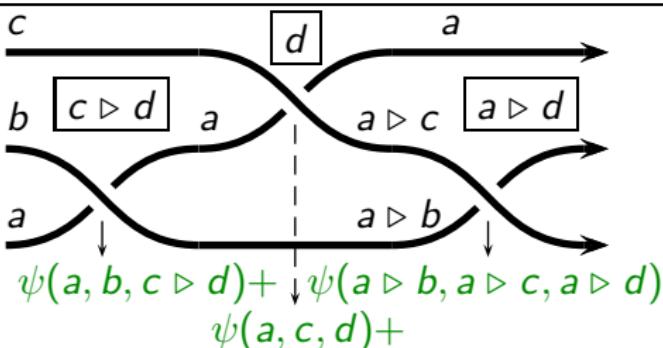
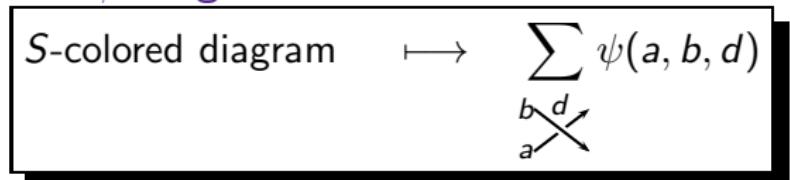
Fix a 3-cocycle $\psi : S \times S \times S \rightarrow \mathbb{Z}$:

$$\begin{aligned} \psi(a, b, c \triangleright d) + \psi(a, c, d) + \psi(a \triangleright b, a \triangleright c, a \triangleright d) = \\ \psi(b, c, d) + \psi(a, b \triangleright c, b \triangleright d) + \psi(a, b, d) \end{aligned}$$

Shadow colorings:



ψ -weight:

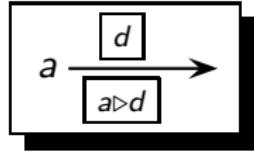


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Shadow colorings:



ψ -weight:

$$S\text{-colored diagram} \longrightarrow \sum_{\substack{b \triangleright d \\ a \times}} \psi(a, b, d)$$

positive braid invariants

colorings &
 \sim
weights

shelf & 3-cocycle

2- and 3-cocycles for Laver tables

Theorem (Dehornoy-L., '14)

① $Z_{\mathbb{R}}^2(A_n) \simeq \mathbb{Z}^{2^n}$ basis: $\phi_{const}(a, b) = 1$ and coboundaries

$$\phi_{q,n}(a, b) = \begin{cases} 1 & \text{if } q \in Col(b), b \notin Col(a \triangleright b), \\ 0 & \text{otherwise.} \end{cases} \quad 1 \leq q < 2^n$$

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Theorem (L., '14)

① $Z_{\text{R}}^k(A_n) \simeq \mathbb{Z}^{P_k(2^n)}$, $P_k(x) = \frac{x^k + x^{\alpha(k)}}{x + 1}$, $\alpha(k) = \begin{cases} 1 & \text{if } k \text{ is even,} \\ 0 & \text{otherwise.} \end{cases}$

② $H_{\text{R}}^k(A_n) \simeq \mathbb{Z}$ for all k .

2- and 3-cocycles for Laver tables

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Remark: 2-cocycles capture the combinatorics of the A_n .

2- and 3-cocycles for Laver tables

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Remark: 2-cocycles capture the combinatorics of the A_n .

Periods via cocycles

$$\pi_n(p) = \min \{ q \mid \phi_{2^{n-1}, n}(p, q) = 1 \}, \quad p < 2^n.$$

$\phi_{1,3}$	1	$\phi_{2,3}$	1 2 3 4 5 6 7 8	$\phi_{3,3}$	1 2 3 4 5 6 7 8	$\phi_{4,3}$	1 2 3 4 5 6 7 8
1	1	1	· 1 · · · · ·	1	1 · 1 · 1 · · ·	1	· · · 1 · · · ·
2	1	2	1 1 · · 1 · · ·	2	· · 1 · · · · ·	2	· · · 1 · · · ·
3	1	3	1 1 · · 1 · · ·	3	1 · 1 · 1 · · ·	3	· 1 · 1 · 1 · ·
4	1	4	· 1 · · · · · ·	4	· · 1 · · · · ·	4	· · · 1 · · · ·
5	1	5	1 1 · · 1 · · ·	5	1 · 1 · 1 · · ·	5	· 1 · 1 · 1 · ·
6	1	6	1 1 · · 1 · · ·	6	1 · 1 · 1 · · ·	6	· 1 · 1 · 1 · ·
7	1	7	1 1 · · 1 · · ·	7	1 · 1 · 1 · · ·	7	1 1 1 1 1 1 1 ·
8	·	8	· · · · · · · ·	8	· · · · · · · ·	8	· · · · · · · ·

$\phi_{5,3}$	1 2 3 4 5 6 7 8	$\phi_{6,3}$	1 2 3 4 5 6 7 8	$\phi_{7,3}$	1 2 3 4 5 6 7 8
1	1 · · · 1 · · ·	1	· 1 · · · 1 · ·	1	1 · 1 · 1 · 1 ·
2	1 · · · 1 · · ·	2	· 1 · · · 1 · ·	2	· · · · · · · ·
3	1 · · · 1 · · ·	3	1 1 1 · 1 1 1 ·	3	1 · 1 · 1 · 1 ·
4	· · · · · · · ·	4	· · · · · · · ·	4	· · · · · · · ·
5	1 · · · 1 · · ·	5	· 1 · · · 1 · ·	5	1 · 1 · 1 · 1 ·
6	1 · · · 1 · · ·	6	· 1 · · · 1 · ·	6	· · · · · · · ·
7	1 · · · 1 · · ·	7	1 1 1 · 1 1 1 ·	7	1 · 1 · 1 · 1 ·
8	· · · · · · · ·	8	· · · · · · · ·	8	· · · · · · · ·

$\phi_{1,3}$	1	$\phi_{2,3}$	1 2 3 4 5 6 7 8	$\phi_{3,3}$	1 2 3 4 5 6 7 8	$\phi_{4,3}$	1 2 3 4 5 6 7 8
1	1	1	· 1 · · · · ·	1	1 · 1 · 1 · · ·	1	· · · · 1 · · ·
2	1	2	1 1 · · 1 · · ·	2	· · 1 · · · · ·	2	· · · · 1 · · ·
3	1	3	1 1 · · 1 · · ·	3	1 · 1 · 1 · · ·	3	· 1 · 1 · 1 · ·
4	1	4	· 1 · · · · ·	4	· · 1 · · · · ·	4	· · · · 1 · · ·
5	1	5	1 1 · · 1 · · ·	5	1 · 1 · 1 · · ·	5	· 1 · 1 · 1 · ·
6	1	6	1 1 · · 1 · · ·	6	1 · 1 · 1 · · ·	6	· 1 · 1 · 1 · ·
7	1	7	1 1 · · 1 · · ·	7	1 · 1 · 1 · · ·	7	1 1 1 1 1 1 1 ·
8	·	8	· · · · · · ·	8	· · · · · · ·	8	· · · · · · ·

$\phi_{5,3}$	1 2 3 4 5 6 7 8	$\phi_{6,3}$	1 2 3 4 5 6 7 8	$\phi_{7,3}$	1 2 3 4 5 6 7 8
1	1 · · · 1 · · ·	1	· 1 · · · 1 · ·	1	1 · 1 · 1 · 1 ·
2	1 · · · 1 · · ·	2	· 1 · · · 1 · ·	2	· · · · · · ·
3	1 · · · 1 · · ·	3	1 1 1 · 1 1 1 ·	3	1 · 1 · 1 · 1 ·
4	· · · · · · ·	4	· · · · · · ·	4	· · · · · · ·
5	1 · · · 1 · · ·	5	· 1 · · · 1 · ·	5	1 · 1 · 1 · 1 ·
6	1 · · · 1 · · ·	6	· 1 · · · 1 · ·	6	· · · · · · ·
7	1 · · · 1 · · ·	7	1 1 1 · 1 1 1 ·	7	1 · 1 · 1 · 1 ·
8	· · · · · · ·	8	· · · · · · ·	8	· · · · · · ·

Main theorem: sketch of proof

Theorem (Dehornoy-L., 14)

$Z_{\text{R}}^2(A_n) \simeq \mathbb{Z}^{2^n}$ basis: $\phi_{\text{const}}(a, b) = 1$ and coboundaries

$$\phi_{q,n}(a, b) = \begin{cases} 1 & \text{if } q \in \text{Col}(b), b \notin \text{Col}(a \triangleright b), \\ 0 & \text{otherwise.} \end{cases} \quad 1 \leq q < 2^n$$

2-cocycle:

$$\phi(a, c) + \phi(a \triangleright b, a \triangleright c) = \phi(b, c) + \phi(a, b \triangleright c)$$

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Step 1. 2-cocycle \implies constant on the last column: $\phi(b, 2^n) = \phi(2^n, 2^n)$.

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0	0	0	1	0	0
q					

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q					

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0	0	0	1	0	0
q					

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Step 4. A change of basis.

$\tilde{\phi}_{1,3}$	1 2 3 4 5 6 7 8	$\tilde{\phi}_{4,3}$	1 2 3 4 5 6 7 8	$\tilde{\phi}_{7,3}$	1 2 3 4 5 6 7 8
1	1	1	. -1· 1 · -1· .	1 1 .
2	1	2	. -1· 1 · -1· .	2	. . -1·
3	1	3	-1· -1 1 -1· -1·	3 1 .
4	1	4	. . . 1	4	. . -1·
5	1	5	. . . 1	5 1 .
6	1	6	. . . 1	6	-1· -1· -1· . .
7	1	7	. . . 1	7 1 .
8	8	8

Part 4

*Bonus: right division ordering
for Laver tables*

Right division for Laver tables

Right division relation:

$$a \mid_r b \iff b = c \triangleright a \text{ for some } c$$

Right division for Laver tables

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Theorem (Dehornoy-L., 14)

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Right division for Laver tables

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Theorem (Dehornoy-L., 14)

- ① \mid_r is a **partial ordering** for A_n .
- ② $a \mid_r b \iff \text{Col}(a) \supseteq \text{Col}(b)$.

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Properties:

❖ Minimal element: 1, maximal element: 2^n .

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Properties:

- ❖ Minimal element: 1, maximal element: 2^n .
- ❖ Linear ordering for $n \leq 2$, **not linear** for $n \geq 3$.

Right division for Laver tables

Right division relation:

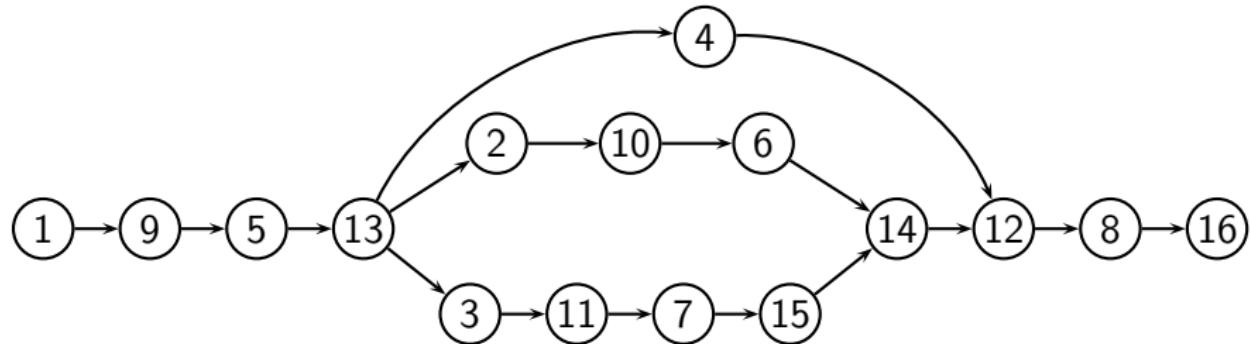
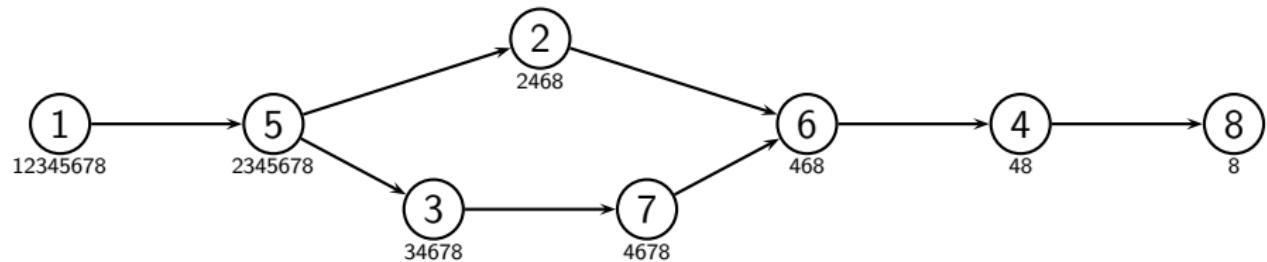
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- ③ $\text{Col}(a) \neq \text{Col}(b)$ for $a \neq b$.

Properties:

- ❖ Minimal element: 1, maximal element: 2^n .
- ❖ Linear ordering for $n \leq 2$, **not linear** for $n \geq 3$.
- ❖ Lattice ordering for $n \leq 4$, **not lattice** for $n \geq 5$.



Main theorem 2: sketch of proof

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Step 7. Anti-symmetry of $|_r$.

A good basis for 2-cocycles

Theorem (Dehornoy-L., 14)

$Z_{\mathbb{R}}^2(A_n) \simeq \mathbb{Z}^{2^n}$ basis: $\phi_{const}(a, b) = 1$ and coboundaries

$$\phi_{q,n}(a, b) = \begin{cases} 1 & \text{if } q \in Col(b), b \notin Col(a \triangleright b), \\ 0 & \text{otherwise.} \end{cases} \quad 1 \leq q < 2^n$$

We saw: ϕ_{const} and 2-coboundaries

$\tilde{\phi}_{q,n} = -d_{\mathbb{R}}^2(\delta_{q,\bullet}) : (a, b) \mapsto \delta_{b,q} - \delta_{a \triangleright b, q} \in \{0, \pm 1\}$, $1 \leq q < 2^n$,

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Change of basis:

$$\phi_{q,n} = \sum_{s|_r q} \tilde{\phi}_{s,n}$$

Digression: Laver tables and branched braids

Theorem (Laver, Drápal, 95)

Operation $p \circ q = p \triangleright (q + 1) - 1$ satisfies

$$(a \circ b) \triangleright c = a \triangleright (b \triangleright c), \quad (a \circ b) \circ c = a \circ (b \circ c),$$

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A_3, \circ	1	2	3	4	5	6	7	8
1	3	5	7	1	3	5	7	1
2	3	6	7	2	3	6	7	2
3	7	3	7	3	7	3	7	3
4	5	6	7	4	5	6	7	4
5	7	5	7	5	7	5	7	5
6	7	6	7	6	7	6	7	6
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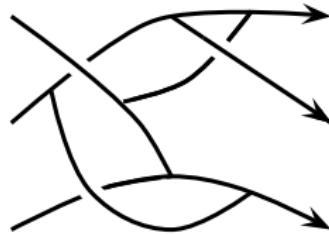
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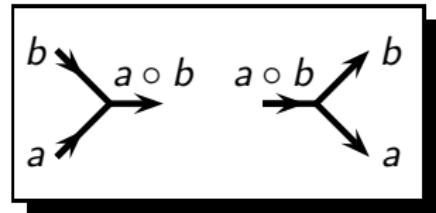
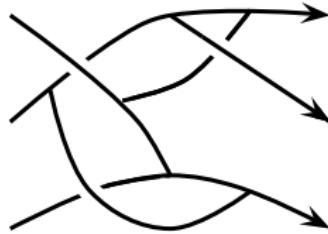
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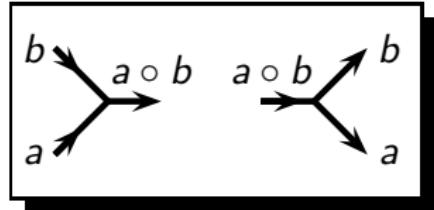
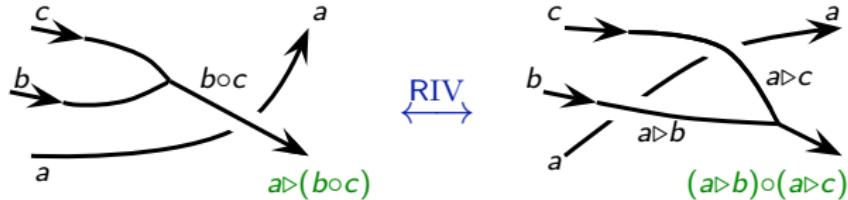
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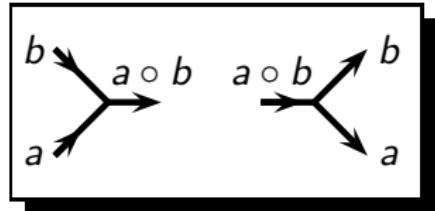
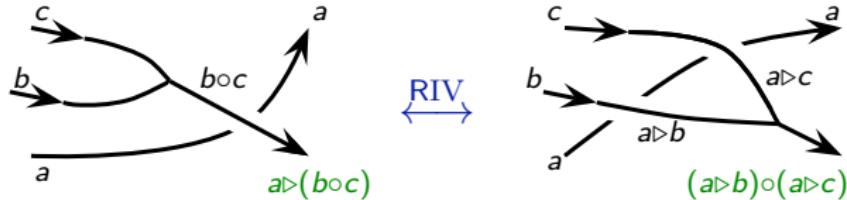
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positive branched braid invariants $\overset{\text{colorings}}{\leadsto} A_n$

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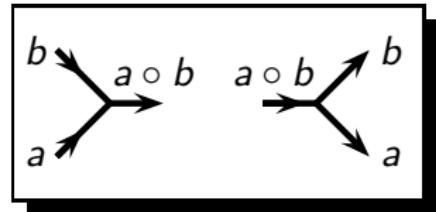
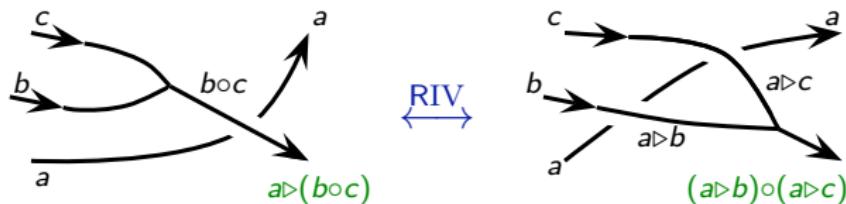
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⚠ Does not work for \mathcal{F}_1 !

Division relations for shelves

	$a \mid_r b$ if $b = c \triangleright a$	$a \mid_l b$ if $b = a \triangleright c$
A_n	is a partial ordering \leadsto a good basis for 2-cocycles	
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 - $b = \gamma \triangleright (\gamma \triangleright \gamma)$, $d(b) = 3$,
 - $a = ((\gamma \triangleright \gamma) \triangleright ((\gamma \triangleright \gamma) \triangleright \gamma)) \triangleright \gamma$, $d(a) = 2$.

Suppose $a \mid_r b$ in \mathcal{F}_1 .

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Suppose $a \mid_r b$ in \mathcal{F}_1 . Then $((1 \triangleright 1) \triangleright ((1 \triangleright 1) \triangleright 1)) \triangleright 1 \mid_r 1 \triangleright (1 \triangleright 1)$ in any A_n .

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- ✿ \mid_r is anti-symmetric, but not transitive on \mathcal{F}_1 .
- ✿ \mid_r strictly sharpens the depth function: $a \mid_r b \Leftrightarrow d(b) = d(a) + 1$.

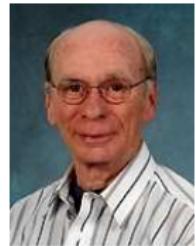
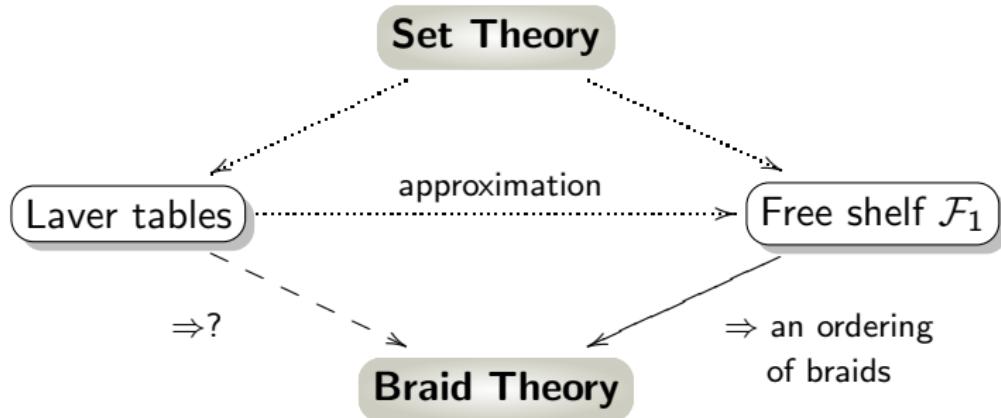
$$b = \gamma \triangleright (\gamma \triangleright \gamma), \quad d(b) = 3,$$

$$a = ((\gamma \triangleright \gamma) \triangleright ((\gamma \triangleright \gamma) \triangleright \gamma)) \triangleright \gamma, \quad d(a) = 2.$$

Suppose $a \mid_r b$ in \mathcal{F}_1 . Then $((1 \triangleright 1) \triangleright ((1 \triangleright 1) \triangleright 1)) \triangleright 1 \mid_r 1 \triangleright (1 \triangleright 1)$ in any A_n . But $8 \nmid_r 4$ in A_3 !

To be continued...

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