

Associative Algebras, Bialgebras and Leibniz Algebras as Braided Objects

Victoria LEBED

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Plan

- 1 Braided Categories and Braided Objects
- 2 “Algebraic” Subcategories of $\mathbf{Br}(\mathcal{C})$
- 3 A Representation Theory for Braided Objects
- 4 A Homology Theory for Braided Objects
- 5 Increasing the Complexity: Multi-Component Braidings
- 6 Braidings as a Unifying Interpretation for Algebraic Structures

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- 5 Increasing the Complexity: Multi-Component Braiding
- 6 Braiding as a Unifying Interpretation for Algebraic Structures

Braided categories

All categories are considered *strict monoidal* in this talk.

Definition

A category \mathcal{C} is called *braided* if it is endowed with a *braiding*, i.e. a family of morphisms $c = \{c_{V,W} : V \otimes W \rightarrow W \otimes V\} \quad \forall V, W \in \text{Ob}(\mathcal{C})$ which is *natural*, i.e. for any $V, W, V', W' \in \text{Ob}(\mathcal{C})$, $f \in \text{Hom}_{\mathcal{C}}(V, V')$, $g \in \text{Hom}_{\mathcal{C}}(W, W')$ one has

$$c_{V',W'} \circ (f \otimes g) = (g \otimes f) \circ c_{V,W}$$

compatible with the tensor product, i.e. $\forall V, W, U \in \text{Ob}(\mathcal{C})$, one has

$$c_{V,W \otimes U} = (\text{Id}_W \otimes c_{V,U}) \circ (c_{V,W} \otimes \text{Id}_U),$$

$$c_{V \otimes W, U} = (c_{V,U} \otimes \text{Id}_W) \circ (\text{Id}_V \otimes c_{W,U}).$$

Braided categories

Definition

A category \mathcal{C} is called *braided* if it is endowed with a *braiding*

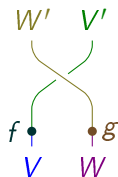
$c = \{c_{V,W} : V \otimes W \rightarrow W \otimes V\} \quad \forall V, W \in \text{Ob}(\mathcal{C})$ which is

✓ *natural*: $c_{V',W'} \circ (f \otimes g) = (g \otimes f) \circ c_{V,W}$

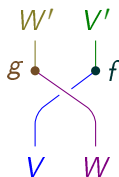
✓ *compatible with \otimes* :

$$c_{V, W \otimes U} = (\text{Id}_W \otimes c_{V,U}) \circ (c_{V,W} \otimes \text{Id}_U),$$

$$c_{V \otimes W, U} = (c_{V,U} \otimes \text{Id}_W) \circ (\text{Id}_V \otimes c_{W,U}).$$



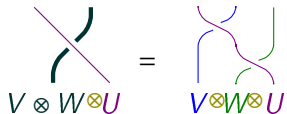
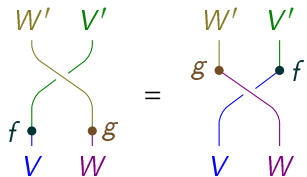
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Braided categories

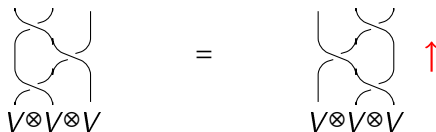


Braided objects

Definition

A *braided object* in \mathcal{C} is an object V endowed with a *braiding*, i.e. a morphism $\sigma_V : V \otimes V \rightarrow V \otimes V$ satisfying (a categorical version of) the *Yang-Baxter equation*:

$$(\sigma_V \otimes \text{Id}_V) \circ (\text{Id}_V \otimes \sigma_V) \circ (\sigma_V \otimes \text{Id}_V) = (\text{Id}_V \otimes \sigma_V) \circ (\sigma_V \otimes \text{Id}_V) \circ (\text{Id}_V \otimes \sigma_V).$$



Yang-Baxter equation \leftrightarrow Reidemeister move III

Braided objects

Definition

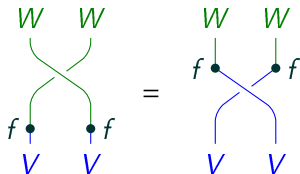
A *braided object* in \mathcal{C} is an object V endowed with a *braiding*, i.e. a morphism $\sigma_V : V \otimes V \rightarrow V \otimes V$ satisfying the *Yang-Baxter equation*:

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A *braided morphism* is a morphism $f : (V, \sigma_V) \rightarrow (W, \sigma_W)$ respecting the braidings:

$$(f \otimes f) \circ \sigma_V = \sigma_W \circ (f \otimes f) : V \otimes V \rightarrow W \otimes W.$$

\rightsquigarrow a category $\mathbf{Br}(\mathcal{C})$.



Braided categories vs. braided objects

braided categories	braided objects
“global” notion	“local” notion

Any object in a braided category is braided.

Remark

We should actually talk about *weakly braided* or *pre-braided* categories / objects, since we do not demand $c_{V,W}$ (or σ_V) to be invertible.

Braided categories vs. braided objects: a digression

Theorem (Folklore)

Denote by \mathcal{C}_{gl-br} the *free braided category* generated by a single object V . Then for each n one has a monoid isomorphism

$$\text{End}_{\mathcal{C}_{gl-br}}(V^{\otimes n}) \xrightarrow{\sim} B_n^+$$

$$\text{Id}_{i-1} \otimes c_{V,V} \otimes \text{Id}_{n-i-1} \mapsto \sigma_i.$$

Here B_n^+ is the *positive Artin braid monoid*:

→ algebraically: generators $\sigma_1, \sigma_2, \dots, \sigma_{n-1}$, subject to relations

$$\sigma_i \sigma_j = \sigma_j \sigma_i \quad \text{if } |i-j| > 1, 1 \leq i, j \leq n-1, \quad (Brc)$$

$$\sigma_i \sigma_{i+1} \sigma_i = \sigma_{i+1} \sigma_i \sigma_{i+1} \quad \forall 1 \leq i \leq n-2; \quad (BryB)$$

→ topologically: braids with positive crossings only.

Braided categories vs. braided objects: a digression

Theorem (Folklore)

Denote by \mathcal{C}_{loc-br} the *free monoidal category* generated by a single *braided object* (V, σ_V) . Then for each n one has a monoid isomorphism

$$\text{End}_{\mathcal{C}_{loc-br}}(V^{\otimes n}) \xrightarrow{\sim} B_n^+$$

$$\text{Id}_{i-1} \otimes \sigma_V \otimes \text{Id}_{n-i-1} \mapsto \sigma_i.$$

Here B_n^+ is the *positive Artin braid monoid*:

→ algebraically: generators $\sigma_1, \sigma_2, \dots, \sigma_{n-1}$, subject to relations

$$\sigma_i \sigma_j = \sigma_j \sigma_i \quad \text{if } |i-j| > 1, 1 \leq i, j \leq n-1, \quad (Brc)$$

$$\sigma_i \sigma_{i+1} \sigma_i = \sigma_{i+1} \sigma_i \sigma_{i+1} \quad \forall 1 \leq i \leq n-2; \quad (BryB)$$

→ topologically: braids with positive crossings only.

Braided categories vs. braided objects: a digression

Theorem (L., 2012)

Denote by $\mathcal{C}_{loc-gl-br}$ the *free symmetric category* generated by a single *braided object* (V, σ_V) . Then for each n one has a monoid isomorphism

$$\text{End}_{\mathcal{C}_{loc-gl-br}}(V^{\otimes n}) \xrightarrow{\sim} VB_n^+$$

$$\text{Id}_{i-1} \otimes c_{V,V} \otimes \text{Id}_{n-i-1} \mapsto \zeta_i,$$

$$\text{Id}_{i-1} \otimes \sigma_V \otimes \text{Id}_{n-i-1} \mapsto \sigma_i.$$

Here VB_n^+ is the *positive virtual braid monoid* (Kauffman, Vershinin):

→ algebraically: generators $\{\sigma_i, \zeta_i, 1 \leq i \leq n-1\}$, subject to

✓ (Br_C) and (Br_{YB}) for the σ_i 's;

✓ (Br_C) and (Br_{YB}) for the ζ_i 's;

$$\zeta_i \zeta_i = 1 \quad \forall i;$$

✓ $\sigma_i \zeta_j = \zeta_j \sigma_i \quad \forall |i-j| > 1,$

✓ $\zeta_i \zeta_{i+1} \sigma_i = \sigma_{i+1} \zeta_i \zeta_{i+1} \quad \forall i$

mixed relations.

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Unital associative algebras

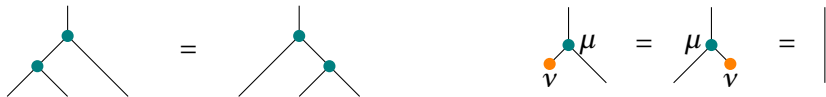
Definition

A *unital associative algebra* (= UAA) in \mathcal{C} is an object V together with morphisms $\mu: V \otimes V \rightarrow V$ and $\nu: I \rightarrow V$, satisfying the associativity and the unit conditions:

$$\mu \circ (\mu \otimes \text{Id}_V) = \mu \circ (\text{Id}_V \otimes \mu) : V^3 \rightarrow V,$$

$$\mu \circ (\nu \otimes \text{Id}_V) = \mu \circ (\text{Id}_V \otimes \nu) = \text{Id}_V.$$

\rightsquigarrow category $\mathbf{Alg}(\mathcal{C})$.



UAAs as braided objects

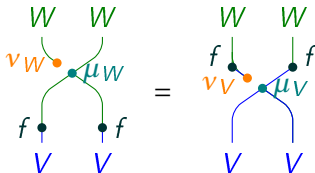
Theorem (L., 2012)

One has a functor

$$\mathbf{Alg}(\mathcal{C}) \longrightarrow \mathbf{Br}(\mathcal{C})$$

$$(V, \mu, \nu) \longmapsto (V, \sigma_{Ass} = \nu \otimes \mu),$$

$$(f : V \rightarrow W) \longmapsto (f : V \rightarrow W).$$



UAAs as braided objects

Definition

Denote by $\mathbf{Br.}(\mathcal{C})$ the category of *pointed braided objects* in \mathcal{C} :

- objects : braided objects V endowed with a "unit" $\nu : I \rightarrow V$;
- morphisms : braided morphisms in \mathcal{C} respecting units, i.e.
 $f \circ \nu_V = \nu_W$ for $f : V \rightarrow W$.

A better theorem (L., 2012)

One has a fully faithful functor

$$\mathbf{Alg}(\mathcal{C}) \hookrightarrow \mathbf{Br.}(\mathcal{C})$$

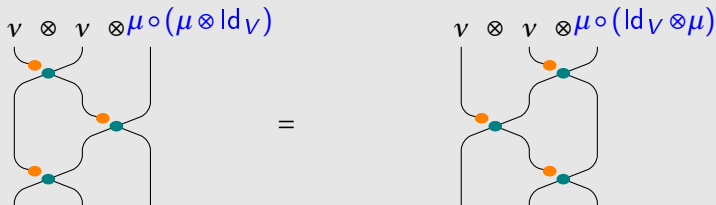
$$(V, \mu, \nu) \longmapsto (V, \sigma_{Ass}, \nu),$$

$$(f : V \rightarrow W) \longmapsto (f : V \rightarrow W).$$

UAAs as braided objects

Remarks

Yang-Baxter for $\sigma_{Ass} \iff$ associativity for μ .



- $\sigma_{Ass} \circ \sigma_{Ass} = \sigma_{Ass} \implies$ highly non-invertible.
- $\sigma = v \otimes \mu + \mu \otimes v - Id$ also encodes the associativity (Nuss, Nichita).
- Dual picture: $\mathbf{coAlg}(\mathcal{C}) \hookrightarrow \mathbf{Br}^*(\mathcal{C})$

$\mathbf{coAlg}(\mathcal{C}) \hookrightarrow \mathbf{Br}^*(\mathcal{C}) \hookrightarrow \mathbf{Alg}(\mathcal{C})$

Unital Leibniz algebras

Definition

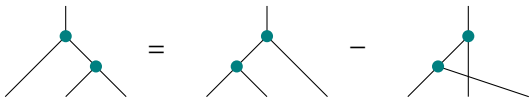
A *unital Leibniz algebra* (= ULA) in a symmetric preadditive category \mathcal{C} is an object V together with morphisms $[,] : V \otimes V \rightarrow V$ and $\nu : I \rightarrow V$, satisfying the Leibniz and the Lie unit conditions:

$$[,] \circ (\text{Id}_V \otimes [,]) = [,] \circ ([,] \otimes \text{Id}_V) - [,] \circ ([,] \otimes \text{Id}_V) \circ (\text{Id}_V \otimes c_{V,V}) : V^{\otimes 3} \rightarrow V,$$

$$[,] \circ (\text{Id}_V \otimes \nu) = [,] \circ (\nu \otimes \text{Id}_V) = 0 : V \rightarrow V.$$

\rightsquigarrow category **Lei**(\mathcal{C}).

A non-commutative version of Lie algebras (Loday, Cuvier).



ULAs as braided objects

Theorem (L., 2012)

One has a fully faithful functor

$$\mathbf{Lei}(\mathcal{C}) \hookrightarrow \mathbf{Br.}(\mathcal{C})$$

$$\begin{aligned} (V, [,], \nu) &\longmapsto (V, \sigma_{Lei} = \nu \otimes [,] + c_{V, V}, \nu), \\ (f : V \rightarrow W) &\longmapsto (f : V \rightarrow W). \end{aligned}$$

ULAs as braided objects

Theorem (L., 2012)

$$\text{Lei}(\mathcal{L}) \hookrightarrow \text{Br.}(\mathcal{L})$$

$$(V, [,], \nu) \mapsto (V, \sigma_{\text{Lei}} = \nu \otimes [,] + c_{V, V}, \nu).$$

Remarks

- Yang-Baxter for $\sigma_{\text{Lei}} \iff$ Leibniz condition for $[,]$.
- A conceptual explication of the choice of the lift of the Jacobi condition.
- σ_{Lei} was previously considered for Lie algebras.
- σ_{Lei} is invertible.
- Dual picture: co-Leibniz algebras.

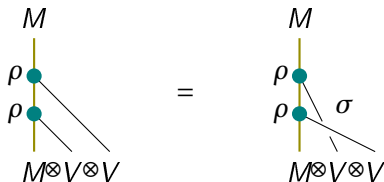
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Braided modules : definition

Definition

A *right braided module* over a braided object (V, σ) in \mathcal{C} is an object $M \in \text{Ob}(\mathcal{C})$ equipped with a morphism $\rho : M \otimes V \rightarrow M$ satisfying

$$\rho \circ (\rho \otimes \text{Id}_V) \circ (\text{Id}_M \otimes \sigma) = \rho \circ (\rho \otimes \text{Id}_V) : M \otimes V \otimes V \rightarrow M.$$



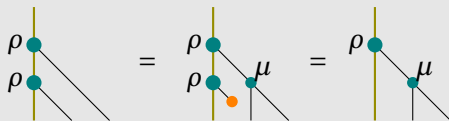
Braided modules : examples

All braided modules are supposed *normalized* here, i.e. $\rho \circ (\text{Id}_M \otimes v) = \text{Id}_M$.

Examples

- UAAs: usual modules over associative algebras

$$\rho \circ (\rho \otimes \text{Id}_V) = \rho \circ (\text{Id}_M \otimes \mu)$$



- ULAs: usual Leibniz modules

$$\rho \circ (\rho \otimes \text{Id}_V) = \rho \circ (\rho \otimes \text{Id}_V) \circ (\text{Id}_M \otimes c_{V,V}) + \rho \circ (\text{Id}_M \otimes [,])$$

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Braided homology

Theorem (L., 2012)

In a preadditive monoidal category \mathcal{C} , take

✓ a braided object (V, σ) ;

✓ a right and a left braided V -modules (M, ρ) and (N, λ) .

Then the morphisms

$${}^{\rho}d_n := \sum_{i=1}^n (-1)^{i-1} {}^{\rho}d_{n;i}, \quad d_n^{\lambda} := \sum_{i=1}^n (-1)^{i-1} d_{n;i}^{\lambda},$$

$${}^{\rho}d_{n;i} := (\rho \otimes \text{Id}_V^{n-1} \otimes \text{Id}_N) \circ (\text{Id}_M \otimes (\sigma_2 \circ \sigma_3 \circ \cdots \circ \sigma_i) \otimes \text{Id}_N)$$

$$d_{n;i}^{\lambda} := (\text{Id}_M \otimes \text{Id}_V^{n-1} \otimes \lambda) \circ (\text{Id}_M \otimes (\sigma_n \circ \cdots \circ \sigma_{i+1}) \otimes \text{Id}_N)$$

define a bidegree -1 tensor bidifferential on $M \otimes V^{\otimes n} \otimes N$.

(Here $\sigma_j = \text{Id}_M \otimes \text{Id}_V^{j-2} \otimes \sigma \otimes \text{Id}_V^{n-j} \otimes \text{Id}_N$.)

Braided homology

Theorem (L., 2012)

$$\rho d_n := \sum_{i=1}^n (-1)^{i-1} \rho d_{n;i}, \quad d_n^\lambda := \sum_{i=1}^n (-1)^{i-1} d_{n;i}^\lambda,$$

$$\rho d_{n;i} := (\rho \otimes \text{Id}_V^{n-1} \otimes \text{Id}_N) \circ (\text{Id}_M \otimes (\sigma_2 \circ \sigma_3 \circ \cdots \circ \sigma_i) \otimes \text{Id}_N)$$

$$d_{n;i}^\lambda := (\text{Id}_M \otimes \text{Id}_V^{n-1} \otimes \lambda) \circ (\text{Id}_M \otimes (\sigma_n \circ \cdots \circ \sigma_{i+1}) \otimes \text{Id}_N)$$

(Here $\sigma_j = \text{Id}_M \otimes \text{Id}_V^{j-2} \otimes \sigma \otimes \text{Id}_V^{n-j} \otimes \text{Id}_N$.)

$$\rho d_{n;i} = \begin{array}{c} \begin{array}{c} | \\ | \\ | \\ | \\ | \\ | \\ | \\ | \\ | \end{array} \\ \begin{array}{c} | \\ | \\ | \\ | \\ | \\ | \\ | \\ | \\ | \end{array} \\ \begin{array}{c} | \\ | \\ | \\ | \\ | \\ | \\ | \\ | \\ | \end{array} \\ \begin{array}{c} | \\ | \\ | \\ | \\ | \\ | \\ | \\ | \\ | \end{array} \\ \begin{array}{c} | \\ | \\ | \\ | \\ | \\ | \\ | \\ | \\ | \end{array} \\ \begin{array}{c} | \\ | \\ | \\ | \\ | \\ | \\ | \\ | \\ | \end{array} \\ \begin{array}{c} | \\ | \\ | \\ | \\ | \\ | \\ | \\ | \\ | \end{array} \end{array}$$

$M \quad \underbrace{\hspace{10em}}_{V^{\otimes n}} \quad N$

Braided homology

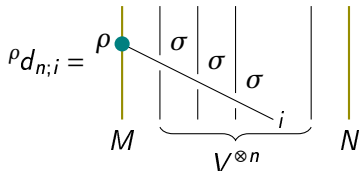
Theorem (L., 2012)

$$\rho d_n := \sum_{i=1}^n (-1)^{i-1} \rho d_{n;i}, \quad d_n^\lambda := \sum_{i=1}^n (-1)^{i-1} d_{n;i}^\lambda,$$

$$\rho d_{n;i} := (\rho \otimes \text{Id}_V^{n-1} \otimes \text{Id}_N) \circ (\text{Id}_M \otimes (\sigma_2 \circ \sigma_3 \circ \cdots \circ \sigma_i) \otimes \text{Id}_N)$$

$$d_{n;i}^\lambda := (\text{Id}_M \otimes \text{Id}_V^{n-1} \otimes \lambda) \circ (\text{Id}_M \otimes (\sigma_n \circ \cdots \circ \sigma_{i+1}) \otimes \text{Id}_N)$$

(Here $\sigma_j = \text{Id}_M \otimes \text{Id}_V^{j-2} \otimes \sigma \otimes \text{Id}_V^{n-j} \otimes \text{Id}_N$.)



Braided homology

Remarks

- A family of differentials \rightsquigarrow linear combinations.
- A *pre-bisimplicial* (or *pre-cubical*) structure on $M \otimes V^{\otimes n} \otimes N$, which can be upgraded into a *weakly bisimplicial* one (a “nice” comultiplication on $V \rightsquigarrow$ degeneracies).
- The construction is *functorial*.
- The differentials can be interpreted in terms of *quantum shuffles* (Rosso).
- A generalization: *Loday's hyperboundaries*

$$M \otimes V^{\otimes n} \otimes N \rightarrow M \otimes V^{\otimes n-k} \otimes N.$$
- Interesting homology morphisms.
- Duality: a cohomology version.

Braided homology: examples

Examples

- UAA V + algebra module $M \rightsquigarrow$
braided object (V, σ_{Ass}) + braided module $M \rightsquigarrow$
bar / Hochschild complex
- ULA V + Leibniz module $M \rightsquigarrow$
braided object (V, σ_{Lei}) + braided module $M \rightsquigarrow$
Leibniz (=non-commutative Chevalley-Eilenberg) complex

algebraic structure \rightsquigarrow chain complex

Braided homology: examples

Examples

- UAA V + algebra module $M \rightsquigarrow$
braided object (V, σ_{Ass}) + braided module $M \rightsquigarrow$
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braided object (V, σ_{Lei}) + braided module $M \rightsquigarrow$
Leibniz (=non-commutative Chevalley-Eilenberg) complex

algebraic structure

case by case
 \rightsquigarrow **braiding** $\xrightarrow{\text{Theorem}}$

chain complex

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Braided systems: definition

Definition

A *braided system* in \mathcal{C} is a family $V_1, V_2, \dots, V_r \in \text{Ob}(\mathcal{C})$ endowed with a *multi-braiding*, i.e. morphisms $\sigma_{i,j} : V_i \otimes V_j \rightarrow V_j \otimes V_i \quad \forall i \leq j$, satisfying the multi-Yang-Baxter equation

$$(\sigma_{j,k} \otimes \text{Id}_i) \circ (\text{Id}_j \otimes \sigma_{i,k}) \circ (\sigma_{i,j} \otimes \text{Id}_k) = (\text{Id}_k \otimes \sigma_{i,j}) \circ (\sigma_{i,k} \otimes \text{Id}_j) \circ (\text{Id}_i \otimes \sigma_{j,k})$$

on all the tensor products $V_i \otimes V_j \otimes V_k$ with $i \leq j \leq k$.

\rightsquigarrow category ${}_r\text{BrSyst}(\mathcal{C})$.



Braided systems: representations and homology

Definition

A *braided system* in \mathcal{C} is a family $V_1, V_2, \dots, V_r \in \text{Ob}(\mathcal{C})$ endowed with a *multi-braiding* $\sigma_{i,j}: V_i \otimes V_j \rightarrow V_j \otimes V_i \ \forall i \leq j$, satisfying YBE on all the tensor products $V_i \otimes V_j \otimes V_k$ with $i \leq j \leq k$.

A *multi-braided module* over a $(\overline{V}, \overline{\sigma}) \in {}_r\mathbf{BrSyst}(\mathcal{C})$ is an $M \in \text{Ob}(\mathcal{C})$ equipped with $(\rho_i: M \otimes V_i \rightarrow M)_{1 \leq i \leq r}$ satisfying $\forall i \leq j$

$$\rho_j \circ (\rho_i \otimes \text{Id}_j) = \rho_i \circ (\rho_j \otimes \text{Id}_i) \circ (\text{Id}_M \otimes \sigma_{i,j}): M \otimes V_i \otimes V_j \rightarrow M.$$

Theorem (L., 2012)

A *bidifferential structure on $M \otimes T(\overline{V})_n^{\rightarrow} \otimes N$* , where $T(\overline{V})_n^{\rightarrow}$ is the direct sum of all the tensor products of type

$$V_1^{\otimes m_1} \otimes V_2^{\otimes m_2} \otimes \dots \otimes V_r^{\otimes m_r}, \quad m_i \geq 0, \sum m_i = n.$$

Example: bialgebras as a braided system

Theorem (L., 2012)

The *groupoid* ${}^*\mathbf{Bialg}(\mathbf{vect}_{\mathbb{k}})$ of bialgebras and bialgebra *isomorphisms* in $\mathbf{vect}_{\mathbb{k}}$ is a subcategory of the *groupoid* of size 2 *bipointed* braided systems:

$${}^*\mathbf{Bialg} \hookrightarrow {}^*_2\mathbf{BrSyst}$$

$$\begin{aligned} (H, \mu, \nu, \Delta, \varepsilon) &\longmapsto \overline{H}_{bi} := (V_1 := H, V_2 := H^*; \quad \nu, \varepsilon^*; \varepsilon, \nu^*; \\ &\quad \sigma_{1,1} := \sigma_{Ass}^r(H), \sigma_{2,2} := \sigma_{Ass}(H^*), \sigma_{1,2} = \sigma_{bi}), \\ f &\longmapsto (f, (f^{-1})^*), \end{aligned}$$

where $\sigma_{bi}(h \otimes l) := \langle l_{(1)}, h_{(2)} \rangle l_{(2)} \otimes h_{(1)}$.

$$\sigma_{H,H} = \begin{array}{c} \diagup \quad \diagdown \\ \mu \quad \nu \\ \diagdown \quad \diagup \end{array}$$

$$\sigma_{H,H^*} = \begin{array}{c} \diagup \quad \diagdown \\ \Delta \quad ev \\ \diagdown \quad \diagup \\ \mu^* \end{array}$$

$$\sigma_{H^*,H^*} = \begin{array}{c} \diagup \quad \diagdown \\ \varepsilon^* \quad \Delta^* \\ \diagdown \quad \diagup \end{array}$$

Example: bialgebras as a braided systems

Theorem (continued)

- *YBE on $H \otimes H \otimes H^*$ \iff the bialgebra compatibility condition.*
- *Invertibility of σ_{bi} \iff H is a Hopf algebra.*
- *Braided modules for $\overline{H}_{bi} \simeq$ right-right Hopf modules over H .*
- *Braided homology for $\overline{H}_{bi} \supseteq$ Gerstenhaber-Schack homology for H .*

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- 6 Braidings as a Unifying Interpretation for Algebraic Structures

Summary: “braided” interpretation for algebraic structures

(multi-)braiding	\leftarrow	algebraic structure
${}_r\mathbf{BrSyst}(\mathcal{C})$	\leftarrow	$\mathbf{Struc}(\mathcal{C})$
YBE	\Leftrightarrow	the defining relation
invertibility	\Leftrightarrow	algebraic properties
braided morphisms	\simeq	structural morphisms
braided modules	\simeq	usual modules
braided differentials	\supseteq	usual differentials

Examples

- \rightarrow UAAs \rightarrow bialgebras \rightarrow self-distributive structures
- \rightarrow ULAs \rightarrow crossed / smash products \rightarrow Yetter-Drinfel'd modules



“algebraic structure = braiding”

Thank you!

